# Explore of the fuzzy community structure integrating the directed line graph and likelihood optimization (Supplemental Materials) 

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(Dated: May 16, 2016)

## I. THE PROOF OF EQ.(3) IN THE MAIN TEXT

In terms of the parameter of $X, \Theta, \Delta, Z$ and $\Omega$, we have:

$$
\begin{equation*}
\Phi=\Theta B_{g} Z D^{-1}, \Psi=\Delta B_{g} Z D^{-1} \tag{1}
\end{equation*}
$$

where $D=\operatorname{diag}(n \Omega)$.
Proof. We have

$$
\begin{equation*}
\phi_{p q}=\sum_{i \in C_{q}} \frac{1}{N_{q}} \theta_{p i} \tag{2}
\end{equation*}
$$

where $i \in C_{q}$ denotes node $i$ is in the cluster $q$ with a size $N_{q}$, and $\frac{1}{N_{q}}$ is the probability of selecting node $i$ from cluster $q$. Further, we have:

$$
\begin{equation*}
\phi_{p q}=\frac{1}{n \omega_{q}} \sum_{i=1}^{n} \theta_{p i}\left(B_{g} Z\right)_{i q} \tag{3}
\end{equation*}
$$

Similarly, we have:

$$
\begin{equation*}
\psi_{p q}=\frac{1}{n \omega_{q}} \sum_{i=1}^{n} \delta_{p i}\left(B_{g} Z\right)_{i q} \tag{4}
\end{equation*}
$$

Thus, we have

$$
\begin{equation*}
\Phi=\Theta B_{g} Z D^{-1}, \Psi=\Delta B_{g} Z D^{-1} \tag{5}
\end{equation*}
$$

The Proof is end.

## II. THE PROOF OF EQ.(5) IN THE MAIN TEXT

$$
\begin{equation*}
L\left(N \mid X, B_{g}\right)=\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \sum_{k=j}^{K} \prod_{j=1}^{n} f\left(\theta_{k j}, a_{i j}\right) f\left(\delta_{k j}, a_{j i}\right) \omega_{k} \tag{6}
\end{equation*}
$$

where $f(x, y)=x^{y}(1-x)^{1-y}$.
Proof. Let $v=i$ denote the event that a node with linkage structure $<$ $a_{i 1}, \ldots, a_{i n}, a_{1 i}, \ldots a_{n i}>$ will be observed in network $N$. Let $y=k$ denote the event that the cluster label assigned to a node is equal to $k$. Let $i \rightarrow_{a_{i j}} j$ denote the event that node $v_{i}$ may or may not link to node $v_{j}$, depending on $a_{i j}$. Let $i \leftarrow a_{j i} j$ denote the event that node $v_{i}$ may or may not be linked by node $v_{j}$ depending on $a_{j i}$. We have:

$$
\begin{align*}
& L\left(N \mid X, B_{g}\right)=\ln \Pi_{i=1}^{n} P(v=i)=\sum_{i=1}^{n} \ln P(v=i) \\
& =\sum_{i=1}^{n} \ln \sum_{k=1}^{K} P(v=i, y=k) \\
& =\sum_{i=1}^{n} \ln \sum_{k=1}^{K} P(v=i \mid y=k) P(y=k) \\
& =\sum_{i=1}^{n} \ln \sum_{k=1}^{K}\left(P\left(<a_{i 1}, \ldots, a_{i n}, a_{1 i}, \ldots a_{n i}>\mid y=k\right) P(y=k)\right) \\
& =\sum_{i=1}^{n} \ln \sum_{k=1}^{K}\left(\prod_{j=1}^{n} P\left(i \rightarrow_{a_{i j}} j \mid y=k\right) P\left(i \leftarrow_{j_{j i}} j \mid y=k\right) P(y=k)\right)  \tag{7}\\
& =\sum_{i=1}^{n} \ln \sum_{k=1}^{K}\left(\prod _ { j = 1 } ^ { n } ( \theta _ { k j } ^ { a _ { i j } } ( 1 - \theta _ { k j } ) ^ { 1 - a _ { i j } } ) \left(\delta_{k j}^{a_{j i}}\left(1-\delta_{k j}\right)^{\left.\left.1-a_{j i}\right) \omega_{k}\right)}\right.\right. \\
& =\sum_{i=1}^{n} \sum_{k=j}^{K} \prod_{j=1}^{n} f\left(\theta_{k j}, a_{i j}\right) f\left(\delta_{k j}, a_{j i}\right) \omega_{k} \\
& =\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \sum_{k=j}^{K} \prod_{j=1}^{n} f\left(\theta_{k j}, a_{i j}\right) f\left(\delta_{k j}, a_{j i}\right) \omega_{k}
\end{align*}
$$

The proof is end.

## III. THE PROOF OF THEOREM 1 IN THE MAIN TEXT

Theorem 1. A local optimum of maximizing Eq.(6) will be guaranteed by recursively calculating Eq.(8) and (9):

$$
\begin{gather*}
\left\{\begin{array}{l}
\theta_{k j}=\frac{\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} a_{i j} \gamma_{l k}}{\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \gamma_{l k}} \\
\delta_{k j}=\frac{\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} a_{j i} \gamma_{l k}}{\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \gamma_{l k}} \\
\omega_{k}=\frac{\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \gamma_{l k}}{n}
\end{array}\right.  \tag{8}\\
\gamma_{l k}=\frac{1}{\sum_{i=1}^{n} b_{i l} \times} \times \\
\sum_{b_{i l} \neq 0} \frac{\prod_{j=1}^{n} f\left(\theta_{k j}, a_{i j}\right) f\left(\delta_{k j}, a_{j i}\right) w_{k}}{\sum_{k=1}^{K} \prod_{j=1}^{n} f\left(\theta_{k j}, a_{i j}\right) f\left(\delta_{k j}, a_{j i}\right) w_{k}}, \tag{9}
\end{gather*}
$$

where $\gamma_{l k}=E\left[z_{l k}\right]=P\left(y=k \mid b=l, X, B_{g}\right)$,i.e., the probability that module $l$ is labeled with community $k$ given $X$ and $B_{g}$.
Proof. Let $L\left(G, Z \mid X, B_{g}\right)$ be the log-likelihood of the joint distribution of line graph $G$ and $Z$ given $X$ and $B_{g}$, we have

$$
\begin{equation*}
L\left(G, Z \mid X, B_{g}\right)=\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \sum_{k=1}^{K} z_{l k}\left(\sum_{j=1}^{n}\left(\ln f\left(\theta_{k j}, a_{i j}\right)+\ln f\left(\delta_{k j}, a_{j i}\right)\right)+\ln \omega_{k}\right) \tag{10}
\end{equation*}
$$

The detailed proof of Eq.(10) are shown in Eq.(21). Considering the expectation of $L\left(G, Z \mid X, B_{g}\right)$ on $Z$, we have:

$$
\begin{equation*}
E\left[L\left(G, Z \mid X, B_{g}\right)\right]=\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \sum_{k=1}^{K} \gamma_{l k}\left(\sum_{j=1}^{n}\left(\ln f\left(\theta_{k j}, a_{i j}\right)+\ln f\left(\delta_{k j}, a_{j i}\right)\right)+\ln \omega_{k}\right) \tag{11}
\end{equation*}
$$

where $E\left[z_{l k}\right]=\gamma_{l k}=P\left(y=k \mid b=l, X, B_{g}\right)$,i.e. the probability that block $l$ will be labeled as cluster $k$ given $X$ and $B_{g}$. Let $J=E\left[L\left(G, Z \mid X, B_{g}\right)\right]+\lambda\left(\sum_{k=1}^{K} w_{k}=1\right)$, we have:

$$
\left\{\begin{array} { l } 
{ \frac { \partial J } { \partial \theta _ { k j } } = 0 } \\
{ \frac { \partial J } { \partial \delta _ { k j } } = 0 } \\
{ \frac { \partial J } { \partial \omega _ { k } } = 0 } \\
{ \frac { \partial J } { \partial \lambda } = 0 }
\end{array} \Rightarrow \left\{\begin{array}{l}
\theta_{k j}=\frac{\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} a_{i j} \gamma_{l k}}{\sum_{l=1}^{L} \sum_{b_{l} \neq 0} \gamma_{l k}} \\
\delta_{k j}=\frac{\sum_{l=1}^{L} \sum_{b_{l i} \neq 0} a_{j i} \gamma_{l k}}{\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \gamma_{l k}} \\
\omega_{k}=\frac{\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \gamma_{l k}}{\sum_{k=1}^{K} \sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \gamma_{l k}}=\frac{\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \gamma_{l k}}{n}
\end{array}\right.\right.
$$

let $P(y=k \mid v=i)$ be the probability that node $i$ belongs to cluster $k$ given $X$ and $B_{g}$, We have:

$$
\begin{equation*}
\gamma_{l k}=P\left(y=k \mid b=l, X, B_{g}\right)=\sum_{b_{i l} \neq 0} \frac{1}{\sum_{i=1}^{n} b_{i l}} P(y=k \mid v=i) \tag{12}
\end{equation*}
$$

where $\frac{1}{\sum_{i=1}^{n} b_{i l}}$ is the probability of selecting node $i$ from block $l$. According to the Bayesian theorem, we have:

$$
\begin{equation*}
P(y=k \mid v=i)=\frac{P(y=k) P(v=i \mid y=k)}{\sum_{k=1}^{K} P(y=k) P(v=i \mid y=k)} . \tag{13}
\end{equation*}
$$

Furthermore, we have:

$$
\begin{equation*}
P(y=k) P(v=i \mid y=k)=\prod_{j=1}^{n} f\left(\theta_{k j}, a_{i j}\right) f\left(\delta_{k j}, a_{j i}\right) \omega_{k} \tag{14}
\end{equation*}
$$

Thus, we have:

$$
\begin{equation*}
\gamma_{l k}=\frac{1}{\sum_{i=1}^{n} b_{i l}} \sum_{b_{i l} \neq 0} \frac{\prod_{j=1}^{n} f\left(\theta_{k j}, a_{i j}\right) f\left(\delta_{k j}, a_{j i}\right) w_{k}}{\sum_{k=1}^{K} \prod_{j=1}^{n} f\left(\theta_{k j}, a_{i j}\right) f\left(\delta_{k j}, a_{j i}\right) \omega_{k}} \tag{15}
\end{equation*}
$$

Moreover, we have:

$$
\begin{align*}
L\left(G \mid X, B_{g}\right) & =\sum_{i=1}^{n} \ln P\left(v=i \mid X, B_{g}\right) \\
& =\sum_{i=1}^{n} \ln \sum_{k=1}^{K} P\left(v=i, y=k \mid X, B_{g}\right) \\
& =\sum_{i=1}^{n} \ln \sum_{k=1}^{K} P\left(y=k \mid v=i, X^{(s)}, B_{g}\right) \frac{P\left(v=i, y=k \mid X, B_{g}\right)}{P\left(y=k \mid v=i, X^{(s)}, B_{g}\right)} \tag{16}
\end{align*}
$$

(by Jensen's inequality)

$$
\begin{aligned}
& \geq \sum_{i=1}^{n} \sum_{k=1}^{K} P\left(y=k \mid v=i, X^{(s)}, B_{g}\right) \ln \frac{P\left(v=i, y=k \mid X, B_{g}\right)}{P\left(y=k \mid v=i, X^{(s)}, B_{g}\right)} \\
& \equiv G\left(X, X^{(s)}\right)
\end{aligned}
$$

Furthermore, we have:

$$
\begin{align*}
G\left(X^{(s)}, X^{(s)}\right) & =\sum_{i=1}^{n} \sum_{k=1}^{K} P\left(y=k \mid v=i, X^{(s)}, B_{g}\right) \ln \frac{P\left(v=i, y=k \mid X, B_{g}\right)}{P\left(y=k \mid v=i, X^{s}, B_{g}\right)} \\
& =\sum_{i=1}^{n} \sum_{k=1}^{K} P\left(y=k \mid v=i, X^{(s)}, B_{g}\right) \ln P\left(v=i \mid X^{(s)}, B_{g}\right)  \tag{17}\\
& =\sum_{i=1}^{n} \ln P\left(v=i \mid X^{(s)}, B_{g}\right) \sum_{k=1}^{K} P\left(y=k \mid v=i, X^{(s)}, B_{g}\right) \\
& =\sum_{i=1}^{n} \ln P\left(v=i \mid X^{(s)}, B_{g}\right) L\left(N \mid X^{(s)}, B_{g}\right) .
\end{align*}
$$

Let $P\left(y=k \mid b=l, X^{(s)}, B_{g}\right)=\gamma_{i k}^{(s)}$, we have:

$$
\begin{align*}
G\left(X, X^{(s)}\right)= & \sum_{l=1}^{L} \sum_{b_{i} \neq 0} \sum_{k=1}^{K} \gamma_{i k}^{(s)} \ln P\left(v=i, y=k \mid X, B_{g}\right)  \tag{18}\\
& -\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \sum_{k=1}^{K} \gamma_{i k}^{(s)} \ln P\left(y=k \mid v=i, X^{(s)}, B_{g}\right) .
\end{align*}
$$

Thus, we have:

$$
\begin{align*}
\operatorname{argmax} G\left(X, X^{(s)}\right)= & \operatorname{argmax}\left(\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \sum_{k=1}^{K} \gamma_{i k}^{(s)} \ln P\left(v=i, y=k \mid X, B_{g}\right)\right. \\
& \left.-\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \sum_{k=1}^{K} \gamma_{i k}^{(s)} \ln P\left(y=k \mid v=i, X^{(s)}, B_{g}\right)\right) \\
& =\operatorname{argmax}\left(\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \sum_{k=1}^{K}\left(\gamma_{i k}^{(s)} \ln P\left(v=i, y=k \mid X, B_{g}\right)\right)\right)  \tag{19}\\
& =\operatorname{argmax} E\left[L\left(G, Z^{(s)} \mid X, B_{g}\right)\right] \\
& =X^{(s+1)} .
\end{align*}
$$

Recall that, the $\Theta^{(s+1)}, \Delta^{(s+1)}$, and $\Omega^{(s+1)}$ of $X^{(s+1)}$ can be computed in terms of $\gamma_{l k}^{(s)}$. Thus, we have:

$$
\begin{equation*}
G\left(X^{(s+1)}, X^{(s)}\right) \geq G\left(X^{(s)}, X^{(s)}\right)=L\left(G \mid X^{(s)}, B_{g}\right) \tag{20}
\end{equation*}
$$

That is to say, the $X^{(s+1)}$ obtained in the current iteration will be not worse than $X^{(s)}$ obtained in last iteration. Thus, we have the theorem and the proof is end.

## IV. THE CALCULATION OF EQ.(10)

$$
\begin{equation*}
L\left(N, Z \mid X, B_{g}\right)=\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \sum_{k=1}^{K} z_{l k}\left(\sum_{j=1}^{n}\left(\ln f\left(\theta_{k j}, a_{i j}\right)+\ln f\left(\delta_{k j}, a_{j i}\right)\right)+\ln \omega_{k}\right) \tag{21}
\end{equation*}
$$

## Proof :

$$
\begin{align*}
& L\left(N, Z \mid X, B_{g}\right) \\
& =\ln \prod_{i=1}^{n} P(v=i, y=y(i)) \\
& =\sum_{i=1}^{n} \ln \sum_{k=1}^{K} m_{i k} P(v=i, y=k) \\
& =\sum_{i=1}^{n} \ln \sum_{k=1}^{K} m_{i k} P(v=i \mid y=k) P(y=k) \\
& =\sum_{i=1}^{n} \sum_{k=1}^{K} \ln (P(v=i \mid y=k) P(y=k))^{m_{i k}}  \tag{22}\\
& =\sum_{i=1}^{n} \sum_{k=1}^{K} m_{i k} \ln (P(v=i \mid y=k) P(y=k)) \\
& =\sum_{i=1}^{n} \sum_{k=1}^{K} m_{i k} \ln \left(\prod_{j=1}^{n}\left(\theta_{k j}^{a_{i j}}\left(1-\theta_{k j}\right)^{1-a_{i j}} \delta_{k j}^{a_{j i}}\left(1-\delta_{k j}\right)^{1-a_{j i}}\right) \omega_{k}\right. \\
& =\sum_{i=1}^{n} \sum_{k=1}^{K} m_{i k}\left(\sum_{j=1}^{n}\left(\ln f\left(\theta_{k j}, a_{i j}\right)+\ln f\left(\delta_{k j}, a_{j i}\right)\right)+\ln \omega_{k}\right) \\
& =\sum_{l=1}^{L} \sum_{b_{i l} \neq 0} \sum_{k=1}^{K} z_{l k}\left(\sum_{j=1}^{n}\left(\ln f\left(\theta_{k j}, a_{i j}\right)+\ln f\left(\delta_{k j}, a_{j i}\right)\right)+\ln \omega_{k}\right)
\end{align*}
$$

The proof is end.

## V. THE CALCULATION OF EQ.(10) IN THE MAIN TEXT

Now, we will discuss how to approximately estimate the prior $P\left(X \mid B_{g}\right)$ based on the information theory. Note that $1 \leq K \leq L=n / g$, which implies, that the coarser the granularity is, the smaller if $K$. It will be shown in the following that a smaller $K$ will indicate a less complexity of $X$. Thus, we have the following: a coarser granularity prefers simpler models, which can be mathematically written as

$$
\begin{equation*}
P\left(X \mid B_{g}\right)=\eta(X)^{g} \tag{23}
\end{equation*}
$$

where the function $\eta(X)$ measures the complexity of $X$ in terms of its parameters. In this paper, we set $\eta(X)=P\left(X \mid B_{1}\right)=P(X)$, the prior of X under $g=1$.

According to Shannon and Weaver, $\ln (1 / P(X))$ is the minimum description length of $X$ with a prior $P(X)$ in its model space. Let $\bar{X}$ denote the optimal coding schema for $X$, and let $C(\bar{X})$ be the coding length (or complexity) of $X$ under this schema. We have

$$
\begin{equation*}
-\ln P\left(X \mid B_{g}\right)=-g \ln P(X)=g C(\bar{X}) \tag{24}
\end{equation*}
$$

Now, to estimate the prior $P\left(X \mid B_{g}\right)$ is to design a coding(or compressing) schema as close to $\bar{X}$ as possible.

As discussed before, the main purpose of the model $X$ is to characterize the behavior of network $N$ in terms of node couplings measured together by its five parameters. Next, we will show that all node couplings can be approximately measured by three compressed parameters instead of original five ones. In this way, one hopes to get a more compact coding schema much closer to $\bar{X}$.

First, we can compress four parameters of $X$,i.e., $\Theta, \Delta, Z$, and $\Omega$, into two

$$
\left\{\begin{array}{l}
\Phi=\Theta B_{1} Z D_{-1}=\Theta Z D_{-1} \\
\Psi=\Delta B_{1} Z D_{-1}=\Delta Z D_{-1}
\end{array}\right.
$$

where $D=\operatorname{diag}(n \Omega)$.
Second, parameter $Z$ can be compressed into a map $y$, where $y(i)=k$ if the entry $(i, k)$ of $B_{1} Z=Z$ is equal to one. Now, given $y, \Phi$, and $\Psi$, node coupling $p_{i j}$ and $q_{i j}$ can be measured by

$$
\begin{equation*}
p_{i j}=\phi_{y(i), y(j)} q_{i j}=\psi_{y(i), y(j)} \tag{25}
\end{equation*}
$$

Equation (25) says that all node coupling can be approximately characterized by $y$, $\Phi$, and $\Psi$. Correspondingly, the compressed coding schema of $X$ is

$$
\begin{equation*}
\hat{X}=\left(K, \Phi_{K \times K}, \Psi_{K \times K}, y_{n \times 2}\right) . \tag{26}
\end{equation*}
$$

Now, we compare the complexity of $X$ and $\hat{X}$

$$
\begin{align*}
C(\hat{X}) & =1 \times\left(-\ln \frac{1}{1}\right)+2 K^{2}\left(-\ln \frac{1}{K^{2}}\right)+2 n\left(-\ln \frac{1}{2 n}\right)  \tag{27}\\
& =2 n K \ln K^{2}+2 n \ln 2 n .
\end{align*}
$$

Note that the dimension of $Z$ is $n \times K$ under $B_{1}$, so we have

$$
\begin{align*}
C(X)= & \times\left(-\ln \frac{1}{1}\right)+3 n K\left(-\ln \frac{1}{n K}\right)+K\left(-\ln \frac{1}{K}\right)  \tag{28}\\
= & 2 n K \ln n K+n K \ln n K+K \ln K .
\end{align*}
$$

It is easy to verify $C(\hat{X}) \ll C(X)$ in the case of $K \ll n$. That is to say that the coding schema closest to $\bar{X}$ that we have found is given by 26 . Thus, we have

$$
\begin{equation*}
C(\bar{X}) \approx C(\hat{X})=2 K^{2} \ln K^{2}+2 n \ln 2 n . \tag{29}
\end{equation*}
$$

The proof is end.

