

Limitation and applicability of modularity measures in community detection

Supplementary Materials 3

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Relationship between the modularity measures Q , D , and the weak definition of community

Given a graph $G = (V, E)$, and partitions $P_K = \{G_1, G_2, \dots, G_K\} = \{(V_1, E_1), (V_2, E_2), \dots, (V_K, E_K)\}$, $K = 1, \dots, |V|$, we denote $Q(P), D(P)$ as the values of Q, D on the partition P .

Relationship between D and the weak definition of community

For modularity measure D , we have the following proposition.

Proposition 2 Let us denote $D(P) = \sum_{i=1}^K D_i$. If for $\forall i$, G_i satisfies the weak definition [1], then we have $D(P) > 0$.

Proof. If G_i satisfies the weak definition, then $D_i > 0$ is valid. It is easy to see that $D(P) = \sum_{i=1}^K D_i > 0$ is valid.

We note that the reverse of **Proposition 2** is not correct, i.e., if $D(P) > 0$ then it is not necessary all G_i ($\forall i$) satisfies the weak definition. An example is shown in Figure 1(b) in the main text. Suppose we have a network which is a 15-clique connected with another two nodes by three edges. In our experimental the optimization of D partition the network into two communities (Shown in Figure 1(b) in the main text). The first one is the 15-clique with D_i value 13.8. The second one is the community with two nodes with D_i value -0.5 . So $D(P) = 13.3 > 0$. However we can simply check that the second community does not satisfy the weak definition.

Relationship between Q and the weak definition of community

For modularity measure Q , we have the following proposition.

Proposition 3 Let us denote $Q(P) = \sum_{i=1}^K Q_i$. If for $\forall i$, G_i satisfies the weak definition [1], then there exists a constant B , and we have $Q_i > B \geq 0$, i.e., Q has a positive lower bound.

Proof

(1) In the simplest situation, $P = \{G_1, G_2\}$, we use $S_1 = L(V_1, V_1), S_2 = L(V_2, V_2), S_{12} = L(V_1, V_2)$ to denote the edges in G_1, G_2 , and the edges between them respectively (as shown in figure S1).

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$$\begin{aligned}
Q_1 &= \frac{S_1}{L} - \left(\frac{2S_1 + S_{12}}{2L} \right)^2 \\
4L^2Q_1 &= 4S_1L - (2S_1 + S_{12})^2 \\
&= 4S_1(S_1 + S_2 + S_{12}) - 4S_1^2 - 4S_1S_{12} - S_{12}^2 \\
&= 4S_1S_2 - S_{12}^2,
\end{aligned}$$

where L represents the number of edges in the whole network. By the weak definition of module,

$$\begin{cases} 2S_1 > S_{12} \\ 2S_2 > S_{12}, \end{cases}$$

we have

$$Q_1 = \frac{4S_1S_2 - S_{12}^2}{4L^2} > 0.$$

Similarly we have

$$Q_2 = \frac{4S_1S_2 - S_{12}^2}{4L^2} > 0.$$

Thus let

$$B = \frac{4S_1S_2 - S_{12}^2}{2L^2} \geq 0.$$

We have $Q = Q_1 + Q_2 > B \geq 0$

(2) Accordingly, for $K = 3$, we have

$$Q_1 > \frac{S_{23}(S_{12} + S_{13})}{L^2} > 0$$

$$Q_2 > \frac{S_{13}(S_{12} + S_{23})}{L^2} > 0$$

$$Q_3 > \frac{S_{12}(S_{13} + S_{23})}{L^2} > 0$$

Thus let

$$B = \frac{S_{23}(S_{12} + S_{13}) + S_{13}(S_{12} + S_{23}) + S_{12}(S_{13} + S_{23})}{L^2} > 0.$$

and we have $Q = Q_1 + Q_2 + Q_3 > B > 0$

(3) In a general case,

$$\begin{aligned}
Q_1 &= \frac{S_1}{L} - \left(\frac{2S_1 + \sum_{j=2}^K S_{1j}}{2L} \right)^2 \\
4L^2Q_1 &= 4S_1L - (2S_1 + \sum_{j=2}^K S_{1j})^2 \\
&= 4S_1 \left(\sum_{i=1}^K S_i + \sum_{i=1}^{K-1} \sum_{j=i+1}^K S_{ij} \right) - 4S_1^2 - 4S_1 \sum_{j=2}^K S_{1j} - \left(\sum_{j=2}^K S_{1j} \right)^2 \\
&= 4S_1 \sum_{i=2}^K S_i + 4S_1 \sum_{i=2}^{K-1} \sum_{j=i+1}^K S_{ij} - \left(\sum_{j=2}^K S_{1j} \right)^2 = (*).
\end{aligned}$$

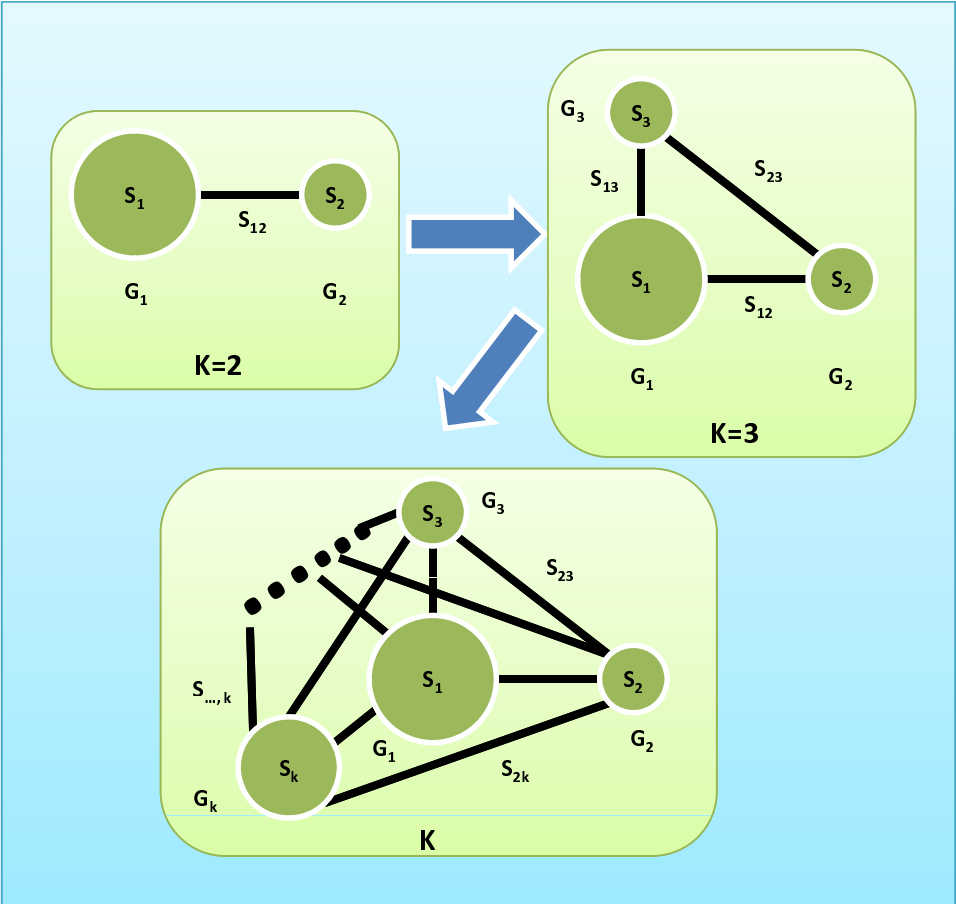


Figure S1: Illustration of different partitions for graph G .

Since all the modules satisfy weak definition, we have

$$\begin{cases} 2S_1 > \sum_{j=2}^K S_{1j} \\ 2S_2 > \sum_{j \neq 2} S_{2j} \\ \dots \\ 2S_K > \sum_{j \neq K} S_{Kj}, \end{cases}$$

so

$$\begin{aligned} (*) &> \left(\sum_{j=2}^K S_{1j} \right) \sum_{i=2}^K \sum_{j \neq i} S_{ij} + 2 \left(\sum_{j=2}^K S_{1j} \right) \sum_{i=2}^{K-1} \sum_{j=i+1}^K S_{ij} - \left(\sum_{j=2}^K S_{1j} \right)^2 \\ &= \left(\sum_{j=2}^K S_{1j} \right) \left(\sum_{i=2}^K \sum_{j \neq i} S_{ij} + 2 \sum_{i=2}^{K-1} \sum_{j=i+1}^K S_{ij} - \sum_{j=2}^K S_{1j} \right) \\ &= \left(\sum_{j=2}^K S_{1j} \right) 4 \sum_{i=2}^{K-1} \sum_{j=i+1}^K S_{ij}, \end{aligned}$$

and further

$$\begin{aligned} Q_1 &> \frac{4 \left(\sum_{j=2}^K S_{1j} \right) \sum_{i=2}^{K-1} \sum_{j=i+1}^K S_{ij}}{4L^2} \\ &= \frac{\left(\sum_{j=2}^K S_{1j} \right) \sum_{i=2}^{K-1} \sum_{j=i+1}^K S_{ij}}{L^2} > 0, \end{aligned}$$

then we can get a lower bound B for Q as

$$B = \sum_{l=1}^K \frac{\left(\sum_{j \neq l} S_{lj} \right) \sum_{i \neq l} \sum_{j=i+1}^K S_{ij}}{L^2}$$

Thus we have $Q = \sum_{l=1}^K Q_l > B > 0$.

Similarly the reverse of **Proposition 3** is not correct, i.e., if $Q(P) > 0$ then it is not necessary all G_i ($\forall i$) satisfies the weak definition. A simple example is given in Figure 1(a) in main text, where there are five 6-cliques, any two of which are connected by eight links. Experimental result shows that the optimization of Q partition the network into five communities and identify every 6-clique as a separate community (Shown in Figure 1(a) in the main text). The optimal value of $Q(P) = 0.280$ is larger than zero. However these five communities all have 15 inner-links and 32 out-links and do not satisfy the weak community definition. We further calculate the lower bound B in **Proposition 3** which has a value 0.32 and is larger than the optimal value of $Q(P)$.

References

- [1] F. Radicchi, C. Castellano, F. Cecconi, V. Loreto, and D. Parisi. Defining and identifying communities in networks. *Proc Natl Acad Sci U S A*, 101(9):2658–2663, March 2004.